

## Exam ALGORITHMS AND DATA STRUCTURES

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NAME: .....

FIRST NAME: .....

MATRICULATION NUMBER: .....

COURSE OF STUDIES: .....

**Note:**

- You have 90 minutes for the exam.
- Please write your name and matriculation number on each sheet.
- Please write clearly. Illegible parts are not corrected and rated as incorrect.
- Cross out concept calculations that should not be counted or make them otherwise identifiable. If several attempts are made to solve a problem, the worst is scored.
- Please use a document-proof pen with blue or black ink and do not use an ink killer or similar. Use only the paper provided.
- Please turn off your electronic devices!

**I declare that I have completed the exam myself and I am aware that the exam will be rated as “failed” if attempted to deceive.**

.....  
(Signature)

Exercise	1	2	3	4	Gesamt
Points	20	20	20	30	90
Scored					

**Exercise 1:**

(a) Let  $f : \mathbb{N} \rightarrow \mathbb{N}$  and  $g : \mathbb{N} \rightarrow \mathbb{N}$  be two functions. State the (mathematical) definition of " $f(n) \in \Omega(g(n))$ ". **(2 Points)**

(b) The function  $T : \mathbb{N} \rightarrow \mathbb{N}$  satisfies  $T(n) = n$  for  $n \leq 100$  and **(4 Points)**

$$T(n) = 64T(n/4) + 5n^3$$

for  $n \geq 101$ . What does the **Master theorem** tell us about the asymptotic behavior of  $T(n)$ ?

(c) Perform an analysis of the Quicksort algorithm. Answer the following questions: **(7 Points)**

- Explain the Divide-step in the QuickSort algorithm.
  
  
  
  
  
  
  
  
  
  
- Explain the Conquer-step in the QuickSort algorithm.
  
  
  
  
  
  
  
  
  
  
- State (without further argument) the time complexity of the Divide-step on an array with  $n$  integers.

- (d) Consider a set of  $n$  integers, each with  $b = 1000$  bits. Explain how **Radix-Sort (7 Points)** with parameter  $r = 10$  sorts these integers. State (without further argument) the resulting time complexity as a function of  $n$ .

**Exercise 2:**

(a) Explain the term **open addressing** in hashing. **(2 Points)**

(b) Explain the term **linear probing** in hashing. **(2 Points)**

(c) Consider a hash table  $T[0 \dots 10]$  with  $m = 11$  entries, and the hash function **(10 Points)**

$$h(k, i) = (2k \bmod 11 + i + 7) \bmod 11.$$

Insert the six keys 11, 46, 25, 47, 24, 13 step by step into the initially empty hash table  $T$ . Use open addressing and linear probing with the given hash function  $h(k, i)$ .

<b>T[0]</b>	<b>T[1]</b>	<b>T[2]</b>	<b>T[3]</b>	<b>T[4]</b>	<b>T[5]</b>	<b>T[6]</b>	<b>T[7]</b>	<b>T[8]</b>	<b>T[9]</b>	<b>T[10]</b>

(d) Explain the problem of **primary clustering** that arises under linear probing. Describe one approach that helps to avoid primary clustering. **(6 Points)**

**Exercise 3:**

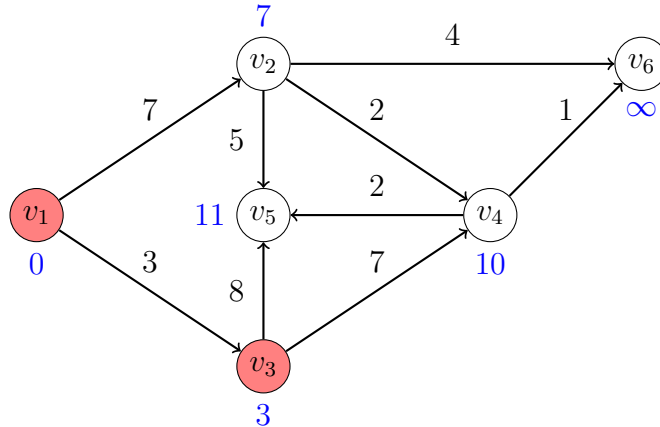
In the **longest common subsequence** problem, we are given a sequence  $X = \langle x_1, x_2, \dots, x_m \rangle$  and a sequence  $Y = \langle y_1, y_2, \dots, y_n \rangle$ . The goal is to find a longest subsequence that occurs in  $X$  as well as in  $Y$ . (Recall that a subsequence of  $X$  results from sequence  $X$  by deleting an arbitrary subset of terms.)

- (a) Explain the **dynamic programming** algorithm for computing a longest common subsequence of  $X$  and  $Y$ . **(8 Points)**

- (b) Perform a time complexity analysis of the dynamic programming algorithm under (a). **(6 Points)**

- State the resulting time complexity as a function of  $m$  and  $n$ .
  
  
  
  
  
  
  
  
  
  
- State the resulting space complexity as a function of  $m$  and  $n$ .

- (c) Construct the dynamic programming table for the sequences  $Y = BCDABA$  **(6 Points)** and  $X = ABDBCAB$ . State a longest common subsequence of  $X$  and  $Y$ .

**Exercise 4:**

The figure above shows an intermediate state of the Dijkstra algorithm, starting at the vertex  $v_1$ . The red vertices are already visited. The blue numbers next to the vertices indicate the current distance to  $v_1$ .

- (a) Insert the missing three lines in the following source code snippet: **(4 Points)**

```

Initialize-Single-Source( $G, s$ );
 $S \leftarrow \emptyset$ ;
 $Q \leftarrow V[G]$ ;
while  $Q \neq \emptyset$  do
    |
    |                                     ;
    |                                     ;
    | for each vertex  $v \in Adj[u]$  do
    | |
    | |                                     ;
    | end
end

```

**Algorithm 1:** Dijkstra algorithm

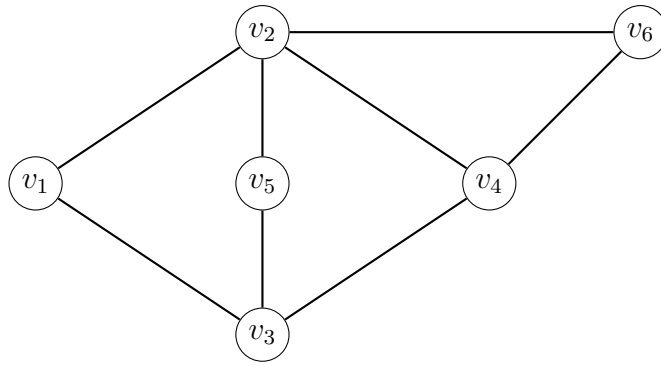
- (b) Consider the next step of the Dijkstra algorithm in the graph presented above: **(6 Points)**

- Which vertex will be selected and handled ?
- State the resulting distances  $d[v]$  for all vertices  $v \in V$ .

- (c) Give the shortest path between  $v_1$  and  $v_6$  at the end of the computation. **(2 Points)**

- (d) Explain why the Dijkstra algorithm can not properly handle edges with negative weight. Give an example that illustrates this problem. **(4 Points)**
- (e) Does the Bellman-Ford algorithm suffer from the same problem with negative edge weights? **(2 Points)**
- (f) What is the running time of the Dijkstra algorithm in terms of  $|V|$  and  $|E|$ ? **(2 Points)**
- (g) Give the proof idea for the time complexity analysis. **(6 Points)**





(h) Draw a Breadth-first search tree, starting at  $v_6$ .

**(4 Points)**